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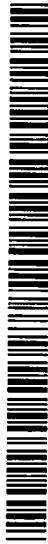
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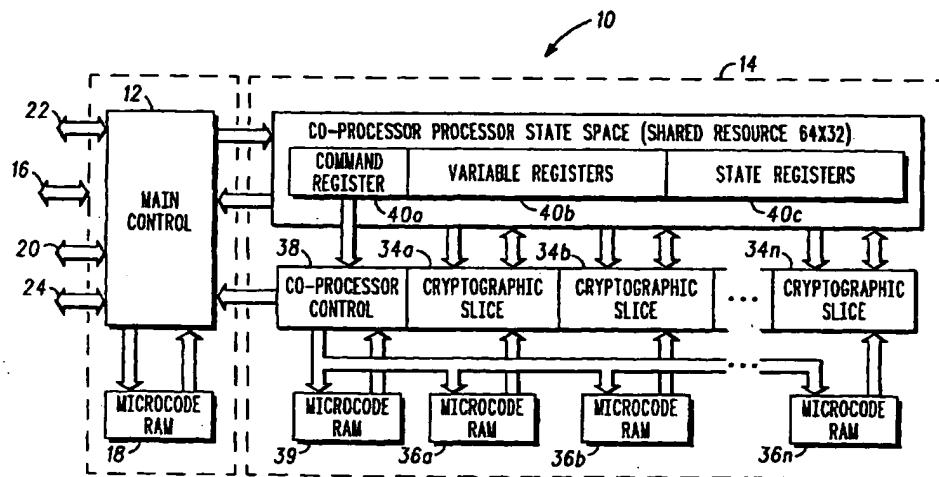
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(54) Title: SCALABLE CRYPTOGRAPHIC ENGINE



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(57) Abstract: A cryptographic engine (10) that includes a scalable cryptographic coprocessor (14) that is controlled by, and separate from, a main engine processor (12). The coprocessor includes a register bank (40) for receiving and storing data packets to be encrypted, and a cryptographic processing device (34a-34n) coupled to the register bank (40) with a processing capacity that is scalable based on application-specific parameters. The coprocessor (14) also includes a control device (38) coupled to the register bank (40) and the cryptographic processing device (34a-34n) for instructing the cryptographic processing device to perform a cryptographic processing operation unique to the cryptographic processing device (34a-34n) based on externally-received processing instructions.

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SCALABLE CRYPTOGRAPHIC ENGINE

Background of the Invention

5 Field of the Invention

The present invention relates generally to data encryption systems, and more particularly to a cryptographic engine that has scalable processing capabilities and that is for use in communications devices requiring secure transmission and
10 reception of data.

Description of Related Art

As the sophistication and importance of broadband communications devices
15 continues to increase in both in commercial and military applications, so does the need to maintain the security of data transmitted by these devices. Currently, such broadband communications devices include cryptographic coprocessors for executing encryption and decryption processing algorithms, while a general-purpose controller typically executes more traditional device processing algorithms.
20 While conventional encryption/decryption engine architectures enable data to be securely transmitted across the bandwidth, such architectures do have associated limitations. For example, in a multi-band multi-mode radio or other similar device requiring high performance cryptographic processing, the underlying cryptographic engine architecture must include several cryptographic co-processors
25 to handle the symmetric encryption and decryption algorithm processing required for secure data transmission and reception. Because the engine footprint must be large enough to accommodate the coprocessors, an engine design including several cryptographic coprocessors tends to increase overall manufacturing costs and limit device design possibilities, as well as increase overall chip power
30 requirements.

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Brief Description of the Drawings

Additional objects and advantages of the present invention will be more readily apparent from the following detailed description of preferred embodiments thereof

5 when taken together with the accompanying drawings in which:

FIG. 1 is a schematic block diagram of a programmable cryptographic engine according to a first preferred embodiment of the present invention;

FIG. 2 is a schematic block diagram of the main processor shown in FIG. 1;

FIG. 3 is a schematic block diagram of one of the coprocessor slices shown in

10 FIG. 1;

FIG. 4 is a diagram of a coprocessor permuter shown in FIG. 3;

FIG. 5 is a diagram of a non-linear function section of the permuter shown in FIG. 3;

FIG. 6 is a diagram of a linear logic unit of the coprocessor shown in FIG. 3;

15 FIG. 7 is a schematic block diagram of the coprocessor controller shown in FIG. 1;

FIG. 8 is an instruction format for the coprocessor microsequencer shown in FIG. 1;

FIG. 9 is a schematic block diagram of a programmable cryptographic engine according to a second preferred embodiment of the present invention;

20 FIG. 10 is a schematic block diagram of a programmable cryptographic engine according to a third preferred embodiment of the present invention;

FIG. 11 is a schematic block diagram of a programmable cryptographic engine according to a fourth preferred embodiment of the present invention; and

25 FIG. 12 is a schematic block diagram of a programmable cryptographic engine according to a fifth preferred embodiment of the present invention.

Detailed Description of a Preferred Embodiment

30 Referring now to the drawings in which like numerals reference like parts, FIG. 1 shows a cryptographic processor engine 10 according to a first embodiment of the present invention. The engine 10, which is typically fabricated on a single semiconductor substrate, includes a main processor 12 for performing general

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purpose engine processing, and a coprocessor 14 having a scalable cryptographic processing capacity that will be described below in more detail. The engine 10 is capable of processing externally-received data packets via any of a variety of techniques including plain-text data encryption, cipher-text data decryption, in-band 5 signaling processing required for bit, word or frame formatting, in-band signal detection, and error detection and correction processing techniques. The engine 10 is capable of executing a single or multiple independent cryptographic algorithms on a large state space, and is capable of supporting both active and shadow techniques, thereby enabling algorithm and key agility through rapid 10 context switching between algorithms and/or state and variable associations.

As shown in FIG. 1, the main processor 12 handles all communications with external devices, such as other communications device processors, through several interfaces. Specifically, the main processor 12 receives configuration and initial state loading instructions through a control interface 16 and performs processing 15 operations, such as asymmetric processing operations used to initialize a data/voice link, based on instructions stored in and fetched from a processor memory 18, such as a microcode RAM, in response to the received configuration and initial state loading instructions. The main processor 12 may also receive data to be processed from an external memory such as a RAM memory (not shown) via 20 a RAM interface 20 if a particular application requires temporary storage of data to be processed, and outputs the data after it has been processed on one or both of the red and black interfaces 22, 24.

Referring to FIG. 2, the main processor 12 includes a 32-bit microsequencer 26 that includes the interfaces 16, 20, 22, 24. Also, the microsequencer 26 is 25 coupled to an arithmetic logic unit (ALU) 28, multiplexers 29a, 29b, a register file 30 and a comparison unit 34. The microsequencer 26 manages all main processor resources based on instructions stored and fetched from the microcode RAM 18.

The ALU 28 is programmed to perform operations such as AND/OR, XOR, ADD and SUBTRACT operations to enable the main processor 12 to perform 30 standard Boolean and arithmetic processing of data packet portions, referred to as data packet elements, selectively output from the multiplexers 29a, 29b. The data packet elements processed by the ALU 28 are output to either the register file 30,

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to the microsequencer 26 for transmission back to external processors through either the red or black interfaces 22, 24, or to the comparison unit 34. The register file 30 functions as a scratchpad memory for temporary data storage when a message received from the microsequencer 26, such as a decode message, 5 requires that the ALU 28 run an algorithm several times, while the comparison unit 34 enables the main processor 12 to compare two 32-bit data packet elements output from the microsequencer 26 and to generate the number of bit errors occurring in one clock cycle for engine processing integrity purposes.

Referring back to FIG. 1, the scalable coprocessor 14 is controlled by the main 10 processor 12 and performs cryptographic processing functions in parallel with the functions performed by the main processor. The coprocessor is operative. The coprocessor is a multiple slice coprocessor with cryptographic processing slices 34a-34n each having corresponding volatile memories, such as microcode RAMs 36a-36n, for storing slice-specific independent control instructions, where the 15 number n is application-specific and is based on design parameters such as chip cost and performance requirements. A coprocessor controller 38 sequentially handles all operations to be performed by the coprocessor based on instructions retrieved from the memory 39 in response to commands received from the main processor 12, and informs the main processor 12 each time a slice has completed a 20 specific processing operation.

Because the slices 34a-34n perform cryptographic processing operations independently from one another and in parallel with the processing of the main processor, the coprocessor 14 is capable of simultaneously performing several independent cryptographic processing operations, with the exact number of 25 possible operations being dependent on the number of slices, at higher, albeit synchronous, clock rates than that of the main processor 12. For example, the coprocessor may receive instructions to loop through a particular encryption algorithm 64 times to generate encrypted data packets, while the main processor need only input the data packets to be encrypted to the coprocessor 14 and 30 subsequently output the resulting encrypted data packets through the red/black interfaces 22, 24. Therefore, the cryptographic processor engine of the present

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invention is capable of supporting data rates equal to the throughput of the coprocessor 14 in a manner that minimizes overall engine power consumption.

As shown in FIG. 1, the coprocessor includes a 64 x 32-bit register bank 40 that includes a command register 40a, variable registers 40b and state registers 40c. The command register 40a, hereinafter referred to as the register R0, is for receiving and loading commands from the main processor 12 and for subsequently sending those commands to the coprocessor controller 38. The variable registers 40b are for storing data packet variables for the slices 34a – 34n such as, for example, a session key for use by the coprocessor 14 in performing cryptographic functions. The state registers 40c are for storing state data such as, for example, channel program data representing a current execution state for an encryption algorithm being executed by one of the slices. The state registers 40c, which are typically accessible by main processors in conventional cryptographic engine architectures, are included as part of the coprocessor 14 in the present invention to enable the number of coprocessor slices to be set according to cryptographic application processing parameters.

FIG. 3 is a block diagram of the cryptographic slice 34a of FIG. 1, with it being understood that the other cryptographic slices 34b - 34n are of like structure and similar function, with the exact function of each slice being dependent on the instructions stored in the microcode RAMs 36a – 36n. The cryptographic slice 34a performs a designated encryption/decryption function for the cryptographic engine 10 that can be independent from processing functions performed by other cryptographic slices that may be included in the engine architecture, even though the slice 34a may share state space in the register bank 40 with other slices. The slice 34a includes an 8-bit microsequencer 50 that provides the slice with processing power and that manages program flow for the slice based on instructions stored in and fetched from the microcode RAM 36a. Because each slice includes a microsequencer such as the microsequencer 50, each slice is capable of performing a processing function independently from other slices. The independent processing capability of each slice produces economies of scale, as the coprocessor can simultaneously execute a number of independent

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cryptographic algorithms corresponding generally to the number of coprocessor slices, thereby increasing overall engine data packet throughput.

The microsequencer 50 controls operation of an input register 54, which in turn provides control bits necessary to move data into and out of the register bank 5 40 and control bits necessary to select various coprocessor cryptographic functions. More specifically, the microsequencer 50 initiates operation of the input register 54 by inputting an address into an input register non-volatile memory 56, such as a microcode RAM, that stores register-specific operating instructions. The microsequencer 50 then latches data output from the RAM 56 to control a permuter 10 58, which has a 160 x 160-bit linear permutation section 60 and a non-linear function unit including a 160 x 144-bit permutation section 62 and a non-linear lookup table section 64 with, for example, sixteen 9:1 lookup tables.

In operation, the data slice 34a is capable of processing ten 32-bit data packet elements input from the register bank 40 to produce four 32-bit encrypted/decrypted 15 data packet elements for return to the state registers 40c every clock cycle. More specifically, five of the input data packet elements, indicated at 70, are input into and permuted by the linear permutation section 60 to create four new 32-bit data packet elements at 72 that are then input into the linear function unit 66. The other five input data packet elements, indicated at 74, are set up by the nonlinear 20 permutation section 62 before being mapped to values in the lookup tables 64 and then output to the linear logic unit 66.

In FIG. 4, the linear permutation section 60 is shown in more detail. The linear permutation section 60 is capable of routing any input bit to any output bit based on control instructions stored in a local section memory (not shown). Each of the five 25 32-bit data packet elements at 70 is input into a 4:1 multiplexer 76. The multiplexer 76 outputs each 32-bit data packet element as four separate 8-bit data packet elements. The 8-bit data packet elements in turn are input into an 8:1 multiplexer 78. The multiplexer 78 outputs each of the 8-bit data packet elements to a multiplexer 80 as eight separate bits of data. The multiplexer 80 sends each of the 30 data bits along with the thirty-one other 1-bit outputs as a permuted 32-bit output to the linear logic unit 66.

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FIG. 5 shows that the non-linear lookup table section 64 is composed of sixteen separate 512 x 1-bit lookup tables, as indicated generally at 81. Sixteen 9-bit data packet elements, resulting from the five 32-bit data packet elements at 74 being permuted by the non-linear permutation section 62, are input into 5 corresponding ones of the lookup tables and mapped to a single bit non-linear table value. The resulting 16-bit data packet element is concatenated at 80 with a previously-generated 16-bit result stored in a delay register 82 to form a fifth new 32-bit data packet element input to the linear function unit at 84.

Still referring to FIG. 5, the non-linear function unit configuration of the present 10 invention enables different slice encryption/decryption algorithms to efficiently utilize the non-linear function unit lookup tables 81. For example, one algorithm may utilize four of the memories as a 4-bit lookup table, while a separate algorithm may simultaneously and independently address eight of the other memories for use as an 8-bit lookup table. Any combination of the lookup tables 81 may be utilized to 15 create a single multi-bit lookup table based on the cryptographic algorithm parameters necessary for a particular slice to perform its stored cryptographic algorithm or function.

Referring now to FIG. 6, the linear function unit 66 is shown in more detail. The linear function unit processes the four 32-bit input data packet elements 72 20 generated by the permutation section 60 and the one 32-bit data packet element 84 generated by the non-linear function unit to produce the four 32-bit linear results at 86. More specifically, the linear function unit 66 either processes the data packet elements In1 – In5 through the EXOR tree indicated generally at 90 or bypasses data packet elements In1 – In4 directly to a multiplexer bank 92 based on 25 instructions from the microsequencer 50. The multiplexer bank 92 then outputs the resulting data packet elements as data packet elements Out1 – Out4 back to the register bank 40 for transmission back to the main processor 12.

Referring to FIG. 7, a block diagram of the coprocessor controller 38 is shown. The coprocessor controller 38 includes an 8-bit microsequencer 94 that provides 30 slice processing power and that manages program flow for a given slice-executed encryption/decryption algorithm in response to commands received from the main processor 12. The microsequencer 94 according to one embodiment of the present

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invention may include a three stage pipeline for performing, for example, fetch, execute and write operations on the program instructions in the memory 39. According to one embodiment of the present invention, the microsequencer 94 does not perform conditional operations, and therefore is capable of operating at 5 high speeds with a simple design.

The coprocessor controller 38 also includes a 4 x 16-bit stack 96 that enables the microsequencer 94 to execute up to four nested loops. The stack 96, in combination with program and loop counting circuitry 98, enables the microsequencer to perform in-line code and loop execution processing for the slices 10, 36a – 36n when the main processor 12 writes initial program count and loop count values to the input register 98 through the command register R0. Consequently, the stack 96 minimizes the number of necessary program execution instructions and thus the size of the slice memories. The stack 96 also enables the engine 10 to run encryption algorithms such as data encryption standard (DES) algorithms 15 that generate 16 different versions of a single session key, and codebook algorithms that necessitate sixteen rounds of actual encryption computation.

The format of the instructions stored in the memory 39 of the microsequencer 94 is shown at 100 in FIG. 8. The 00 Continue Opcode is used to execute in-line code, while the 01 Loop Start Opcode is used to signify the start of a loop. When 20 the microsequencer 94 executes the 01 Loop Start Opcode, a value contained in the Loop cnt field and the next instruction address are pushed onto the stack 96. The 10 Loop End Opcode signifies the end of a loop and causes a current count value stored in the program and loop counting circuitry 98 to be decremented when executed by the microsequencer 94. If the current count is not equal to zero, the 25 program counter component 102 of the program and loop counting circuitry 98 is loaded with a current Loop cnt from the stack 96. If the count is equal to zero, the current entry is removed from the stack 96, and the microsequencer 94 continues to execute in-line instructions.

It should be appreciated at this point that the cryptographic processing engine 30 10 of the present invention is designed so that the configuration of the coprocessor controller 38 remains the same regardless of the number of coprocessor slices. In addition, the number of instructions that must be generated by the coprocessor

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controller 38 to perform equivalent functional operations is actually reduced as the number of cryptographic slices is increased. The decrease in the number of required instructions subsequently enables the size of the corresponding coprocessor controller memory 39 to be reduced.

5 FIGs. 9 – 12 show additional embodiments of the scalable cryptographic coprocessor of the present invention at 110, 110', 110" and 110"", respectively. The cryptographic coprocessor 110 is a single-slice cryptographic coprocessor with a slice 134a. The cryptographic coprocessor 110' is a double-slice cryptographic coprocessor with slices 134a', 134b'. The cryptographic coprocessor 110" is a 10 triple-slice cryptographic coprocessor with slices 134a"-134c". The cryptographic coprocessor 110"" is a quadruple-slice cryptographic coprocessor with slices 134a""-134d"". Of the engines shown, the cryptographic engine 110 has the least amount of processing power, and therefore the lowest associated cost and the smallest footprint, while the cryptographic engine 110"" has the most processing power, and 15 therefore the highest associated cost and largest footprint. Additional slices could be added to the embodiments shown to increase processing power if necessary.

The embodiments shown in FIGs. 9-12 therefore show that the scalable architecture of the present invention provides design flexibility that enables a cryptographic engine to be configured to fit within specific application processing 20 and cost parameters. This design flexibility eliminates the need for multiple cryptographic engines to support high performance cryptographic algorithm processing. The single engine architecture of the present invention also enhances overall cryptographic processing performance when compared to conventional cryptographic engine architectures, as it has a target throughput 25 encryption/decryption processing rate of 50 – 200 Mbps for a combination of algorithms and associations.

The above-described scalable programmable cryptographic engine of the present invention is designed to support high performance communications applications such as personal computer cards, network encryption systems, and 30 satellite communications, and can be embedded in applications such as programmable and handheld radios, avionics equipment, network security systems, telephony and numerous other applications requiring secure data transmission and

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reception capabilities. The engine can ensure data integrity on both personal computers and networks and at the same time maintain interoperability with numerous cryptographic algorithm implementations.

While the above description is of the preferred embodiment of the present
5 invention, it should be appreciated that the invention may be modified, altered, or varied without deviating from the scope and fair meaning of the following claims.

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Claims

What is claimed is:

1. A cryptographic engine, comprising:
 - 5 a main processor for performing only asymmetric processing;
 - a control interface for enabling main processor external communication; and
 - a coprocessor subordinate to and separate from the main processor for executing symmetric cryptographic processing algorithms stored therein, the coprocessor having a processing capacity that is scalable based on application-specific parameters.
2. The cryptographic engine of claim 1, wherein the coprocessor includes at least one processing slice.
- 15 3. The cryptographic engine of claim 1, wherein the coprocessor includes multiple processing slices each for independently executing a corresponding cryptographic processing algorithm.
- 20 4. The cryptographic engine of claim 1, wherein the coprocessor has an aggregate processing throughput in a range of approximately 50 Mbps – 200 Mbps.
5. The cryptographic engine of claim 1, wherein the main processor comprises:
 - 25 a multi-bit microsequencer for managing resources of the main processor;
 - a register for receiving data packet elements to be processed from the microsequencer;
 - an arithmetic logic unit coupled to the register and the microsequencer for receiving the data packet elements from the register, for performing standard Boolean and arithmetic operations on the received data packet elements, and for
 - 30 sending processing status information on the received data packet elements to the microsequencer; and

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5 a comparison unit coupled to the microsequencer, the register and to the arithmetic logic unit for enabling the main processor to cyclically compare multi-bit data packet elements from the microsequencer and to generate and transmit to the microsequencer a subsequent number of bit errors to maintain engine processing integrity.

10 6. The cryptographic engine of claim 1, wherein the coprocessor includes a register bank for holding data packet elements to be processed by the coprocessor, and for holding data packet elements processed by the coprocessor to be returned to the main processor.

15 7. The cryptographic engine of claim 6, wherein the coprocessor further includes:

15 a microsequencer for managing coprocessor program flow based on the data packet elements to be processed in the register bank, and on the processed data packet elements in the register bank to be returned to the main processor; and

20 a coprocessor slice including an input register for moving data in and out of the register bank and for selecting stored cryptographic function codes in response to instructions received from the main processor.

20

25 8. The cryptographic engine of claim 7, wherein the main processor is for controlling operation of the coprocessor by writing to a register R0 in the register bank.

25 9. The cryptographic engine of claim 7, wherein the coprocessor slice is for processing 10 32-bit input data packet elements to produce four 32-bit output data packet elements during a single clock cycle.

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10. The cryptographic engine of claim 9, wherein the coprocessor slice further comprises:

a linear permutation section for permuting a first five of the ten 32-bit input data packet elements to create four new 32-bit permuted data packet elements;

5 a non-linear permutation section for processing a second five of the ten 32-bit input values to create a fifth 32-bit permuted data packet element;

a linear logic unit for processing the five 32-bit permuted data packet elements to generate four 32-bit encrypted data packet elements to be output to the main processor.

10

11. The cryptographic engine of claim 10, wherein the non-linear function unit comprises:

a non-linear permutation section for permuting the second five 32-bit data packet elements into sixteen 9-bit data packet elements;

15 sixteen 512 x 1 lookup tables each for mapping a non-linear variable to one of the 9-bit data packet elements input thereto to generate a 16-bit non-linear data packet element; and

a delay register for causing the 16-bit non-linear data packet to be concatenated with a previously generated 16-bit non-linear data packet to produce

20 the fifth 32-bit permuted data packet input into the linear logic unit.

12. The cryptographic engine of claim 1, wherein the main processor includes a nested loop stack for controlling nested loop operation of the coprocessor in response to receiving nested loop instructions from the main processor.

25

13. A coprocessor for a cryptographic engine that is controlled by, and separate from, a main processor, the coprocessor comprising:

a register bank for receiving and storing data packet elements to be encrypted;

a cryptographic processing device coupled to the register bank and having a

30 processing capacity that is scalable based on application-specific parameters; and

a coprocessor control device coupled to the register bank and the cryptographic processing device for instructing the cryptographic processing device

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to perform a cryptographic processing operation unique to the cryptographic processing device.

14. The coprocessor of claim 13, wherein the cryptographic processing device
5 includes a microsequencer for controlling cryptographic processing device
operation in response to operating instructions received from the coprocessor
control device.

15. The coprocessor of claim 14, wherein the cryptographic processing device
10 includes a permuter for performing both linear and non-linear permutation
processing on data packet elements input from the register bank in response to
control instructions received from the microsequencer.

16. The coprocessor of claim 15, wherein the permuter includes a plurality
15 of non-linear lookup tables each containing non-linear variables to be mapped to
input data and that can be utilized separately or in combination based on
cryptographic algorithm processing requirements.

17. The coprocessor of claim 13, wherein the cryptographic processing device
20 comprises a plurality of cryptographic processing devices each coupled to the
register bank and the coprocessor control device and each for performing a unique
cryptographic processing operation independently from others of the cryptographic
processing devices.

25 18. The coprocessor of claim 17, wherein each of the plurality of
cryptographic processing devices includes a microsequencer for controlling
cryptographic processing device operation independently from operation of other
cryptographic processing device microsequencers and in response to operating
instructions received from the coprocessor control device.

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19. The coprocessor of claim 17, wherein the cryptographic processing device includes a permuter for performing both linear and non-linear permutation processing on data packet elements input from the register bank
- 5 20. The coprocessor of claim 19, wherein the permuter includes a plurality of non-linear lookup tables each containing non-linear variables to be mapped to input data and that can be utilized separately or in combination based on cryptographic algorithm processing requirements.
- 10 21. The coprocessor of claim 13, wherein the coprocessor control device includes a stack for enabling the cryptographic processing device to loop through a corresponding cryptographic processing algorithm a predetermined number of times.

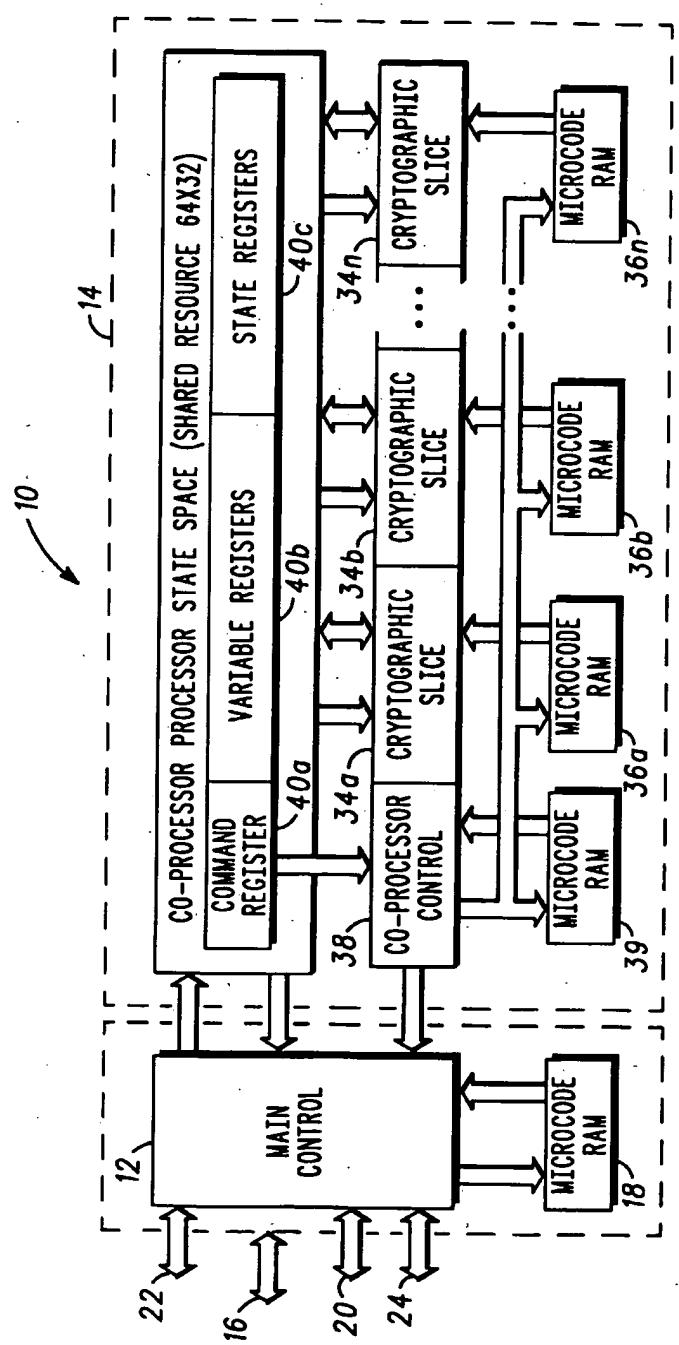


FIG. 1

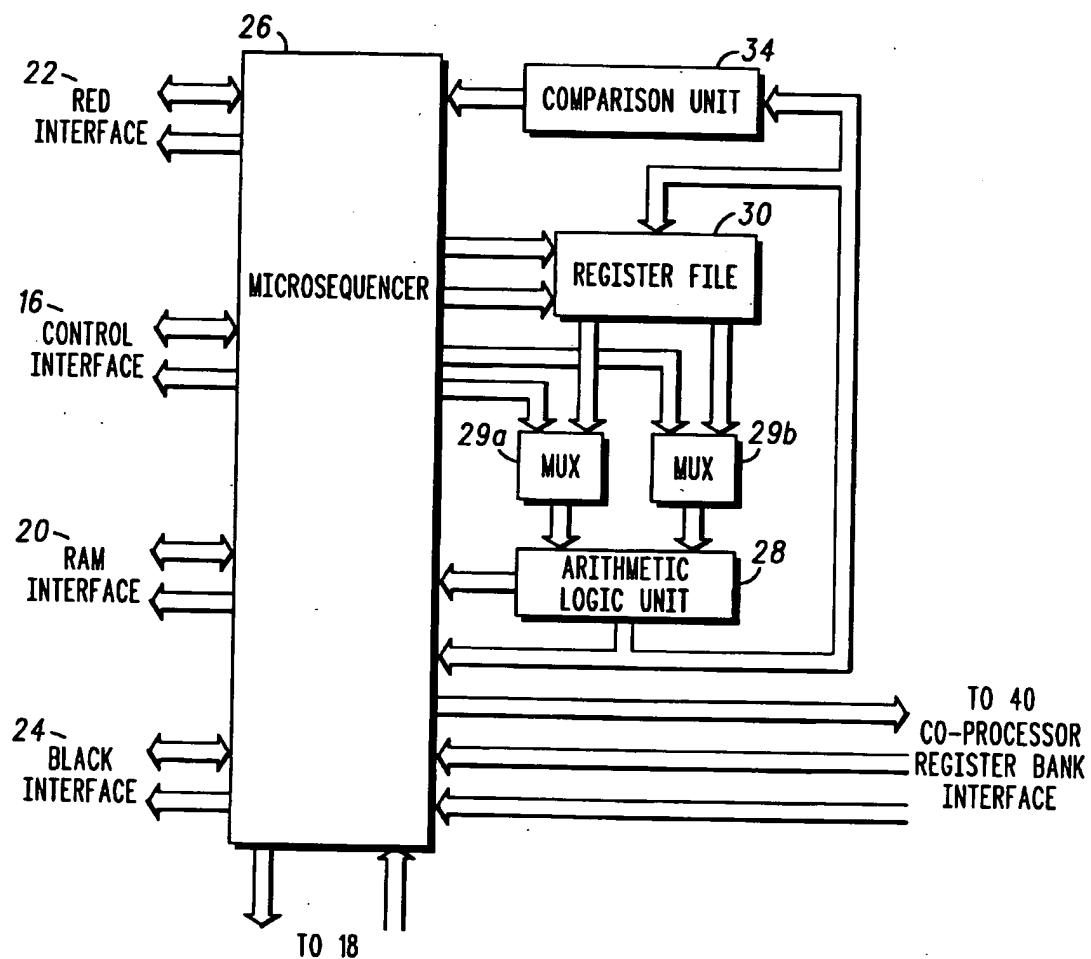
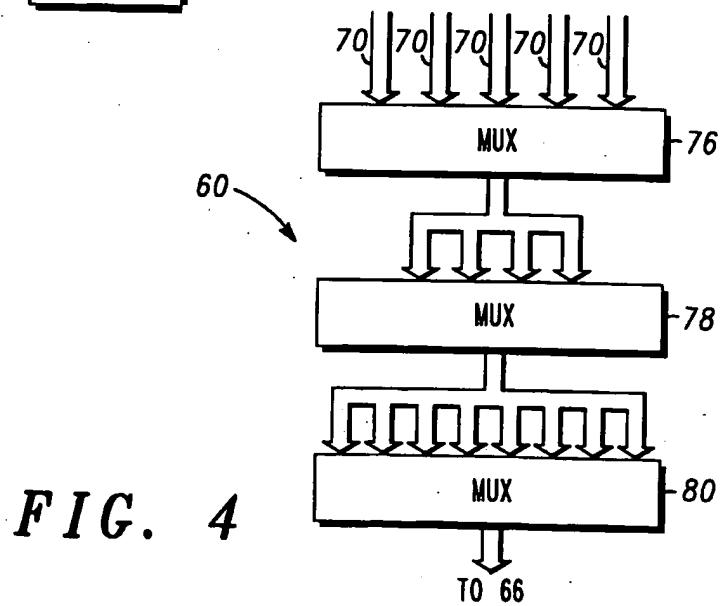
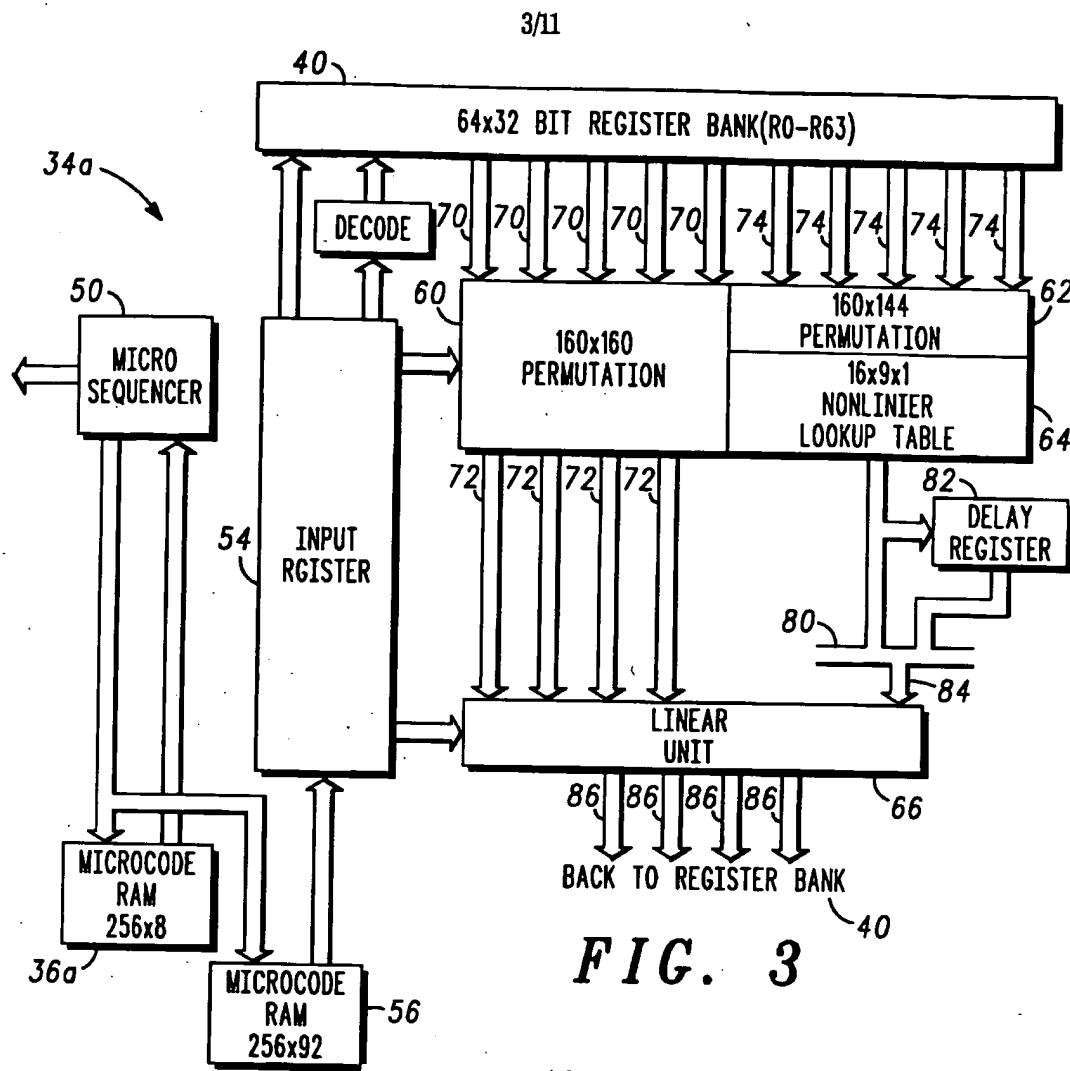


FIG. 2



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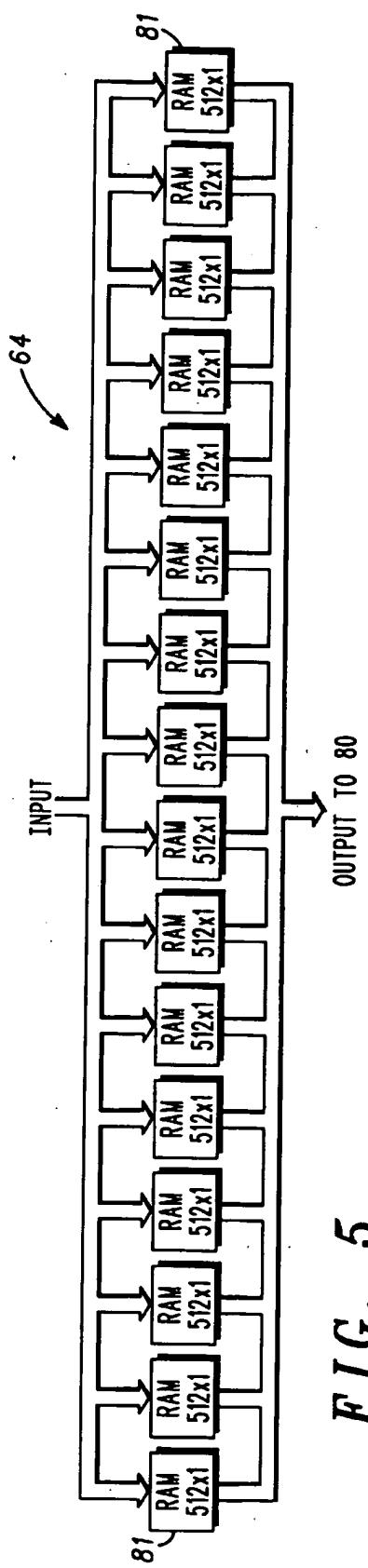


FIG. 5

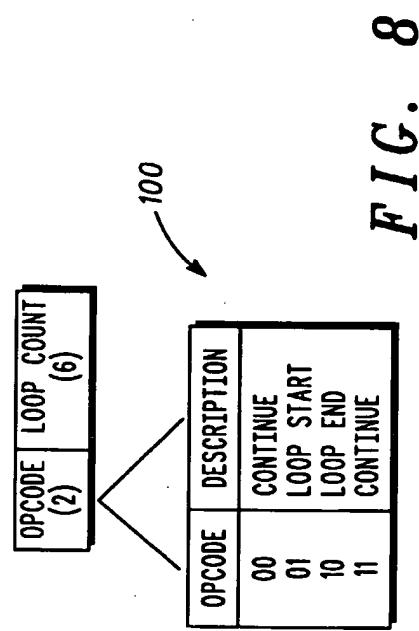


FIG. 8

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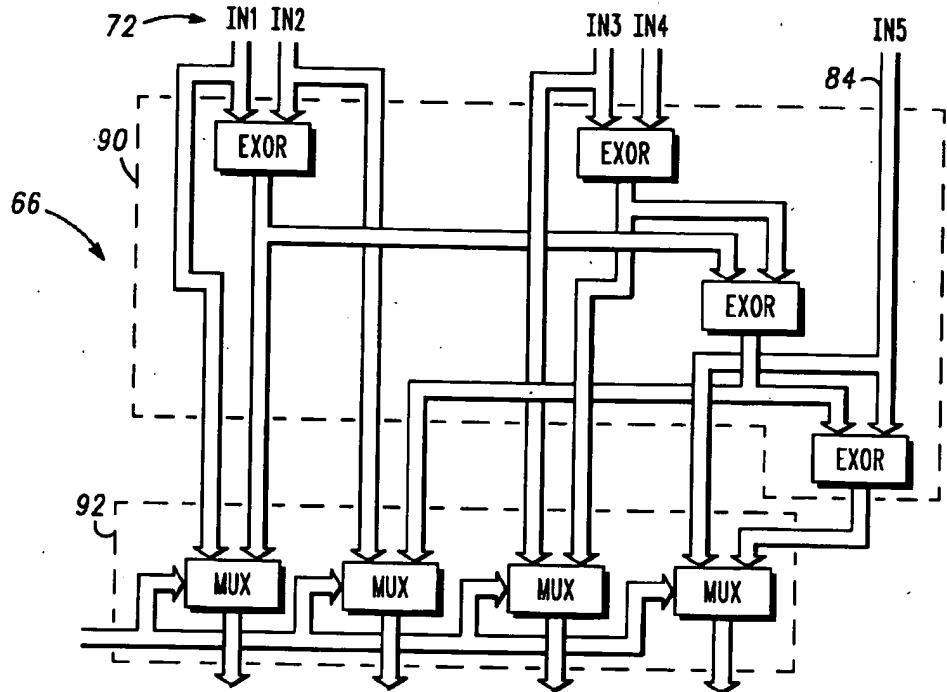
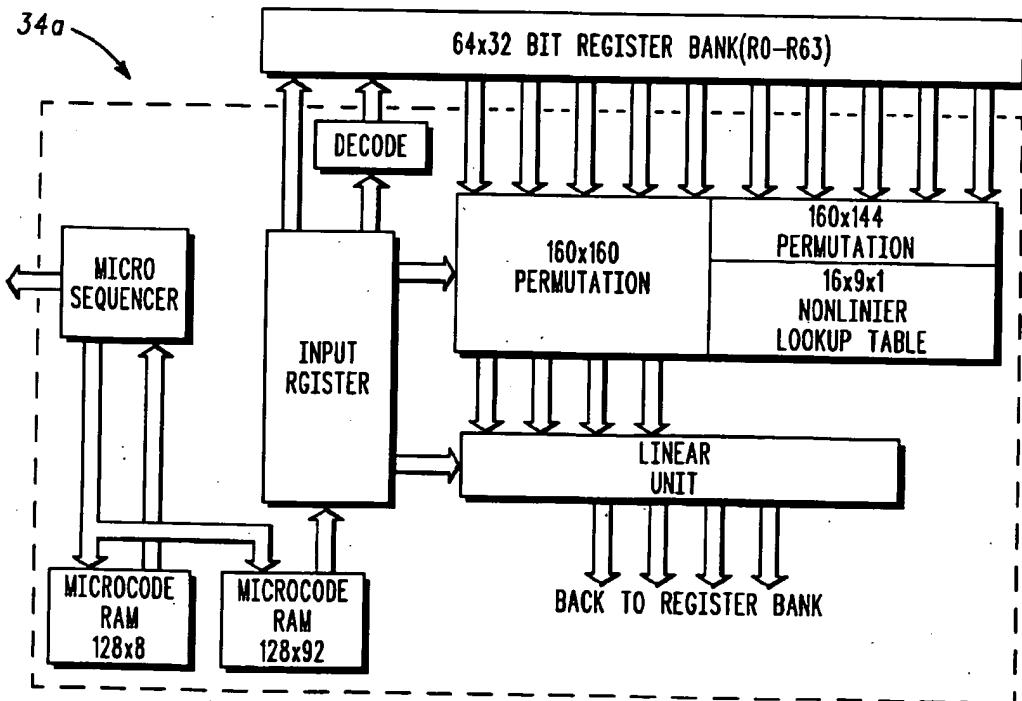
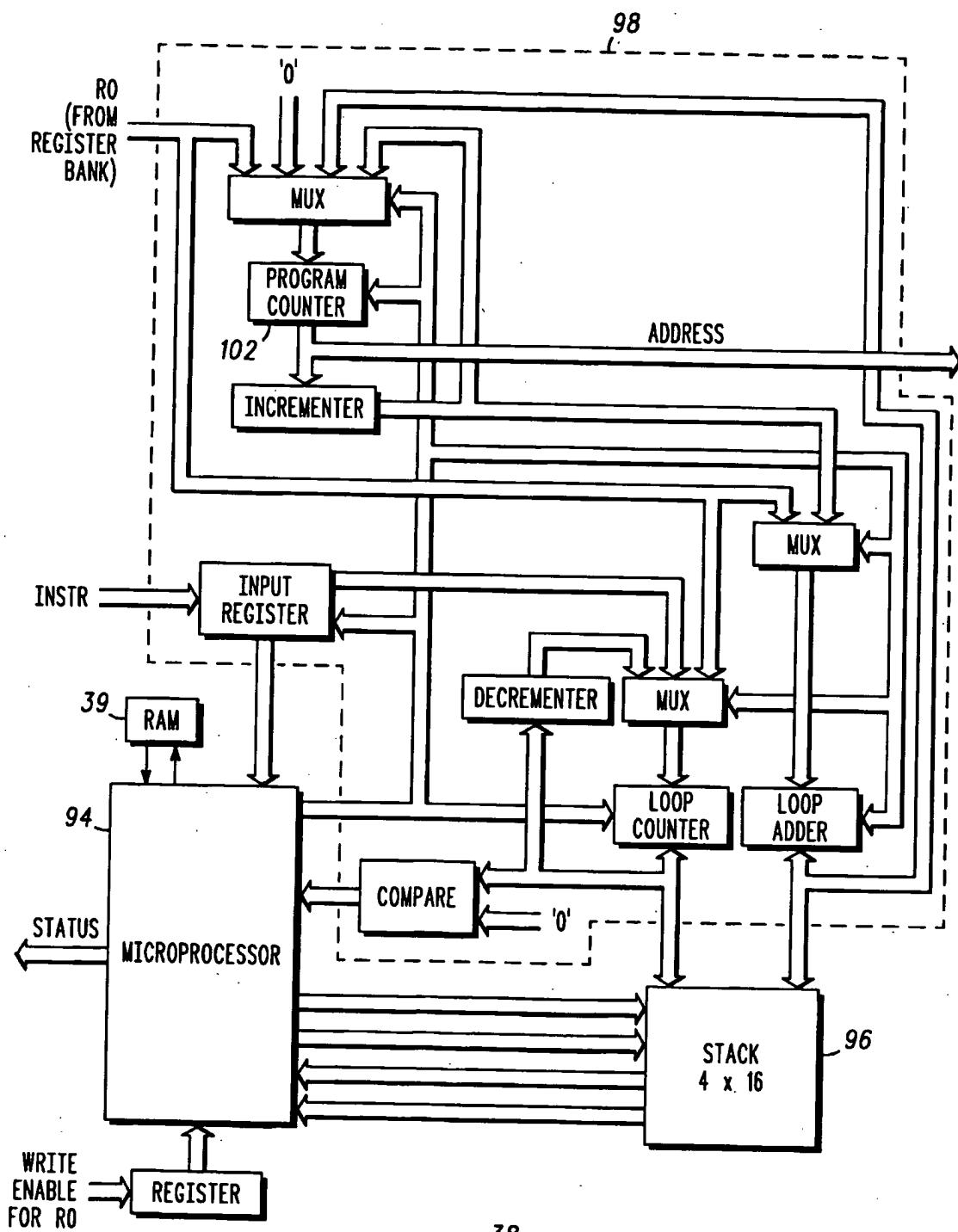


FIG. 6

FIG. 9



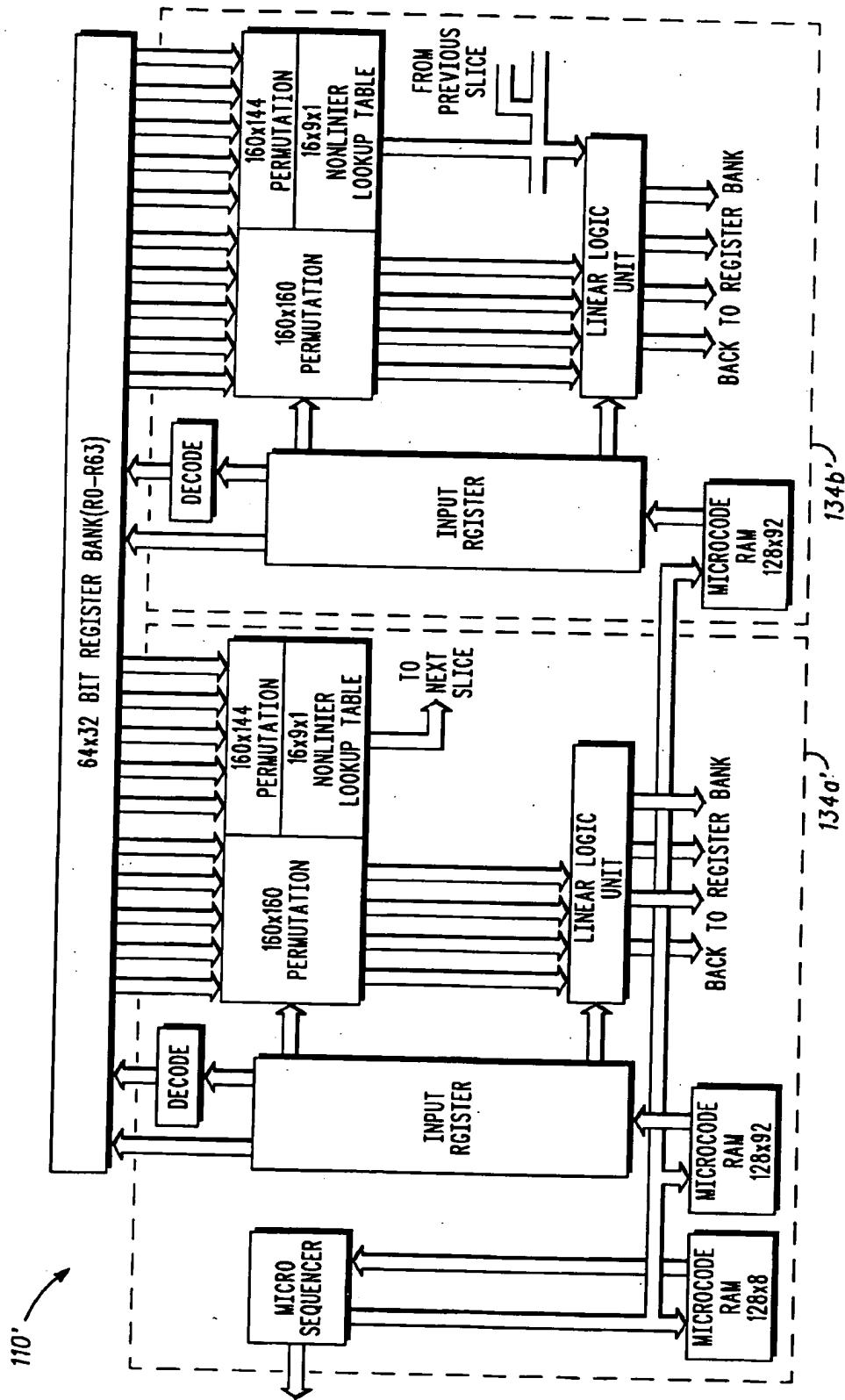
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FIG. 7

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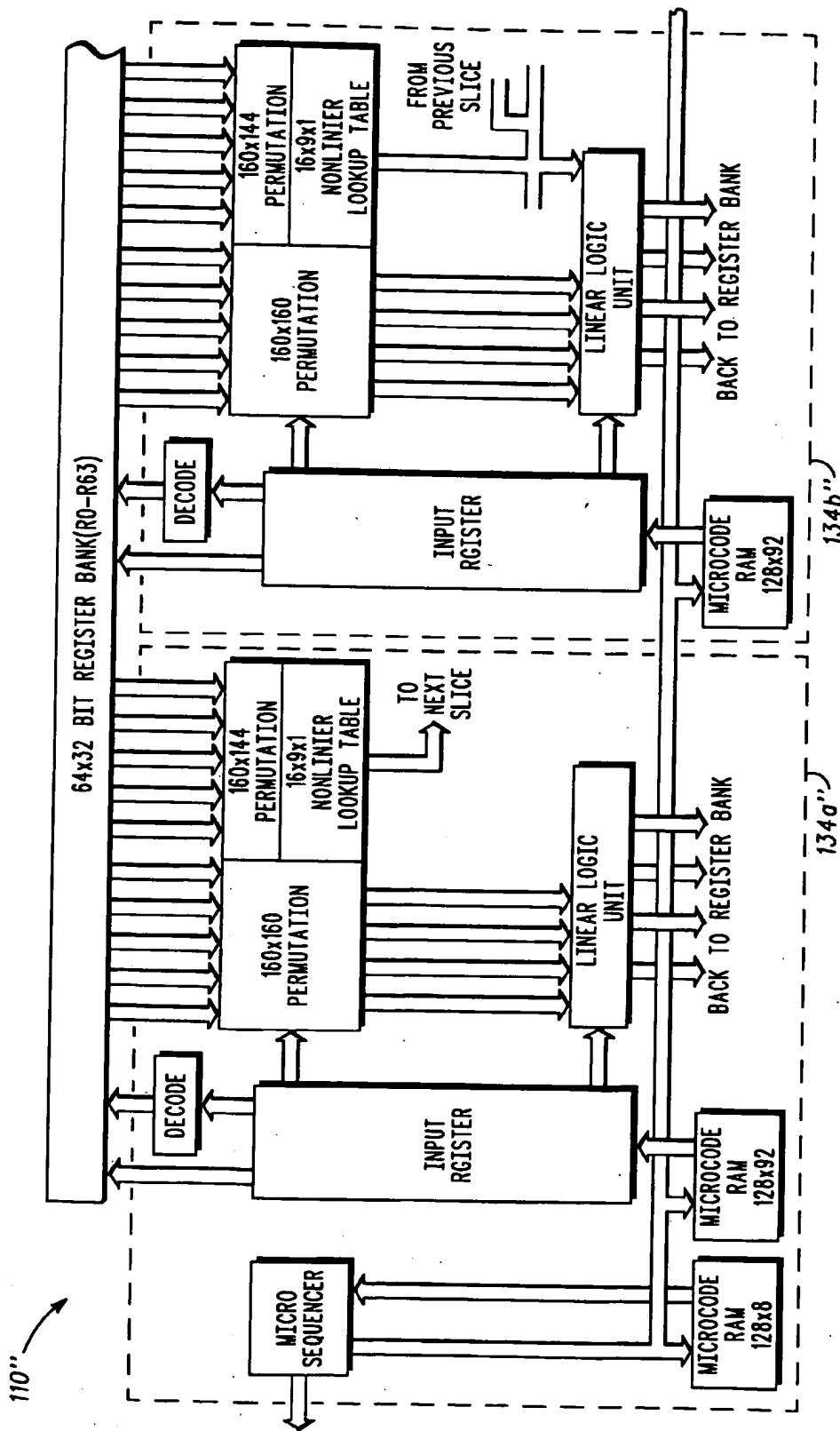
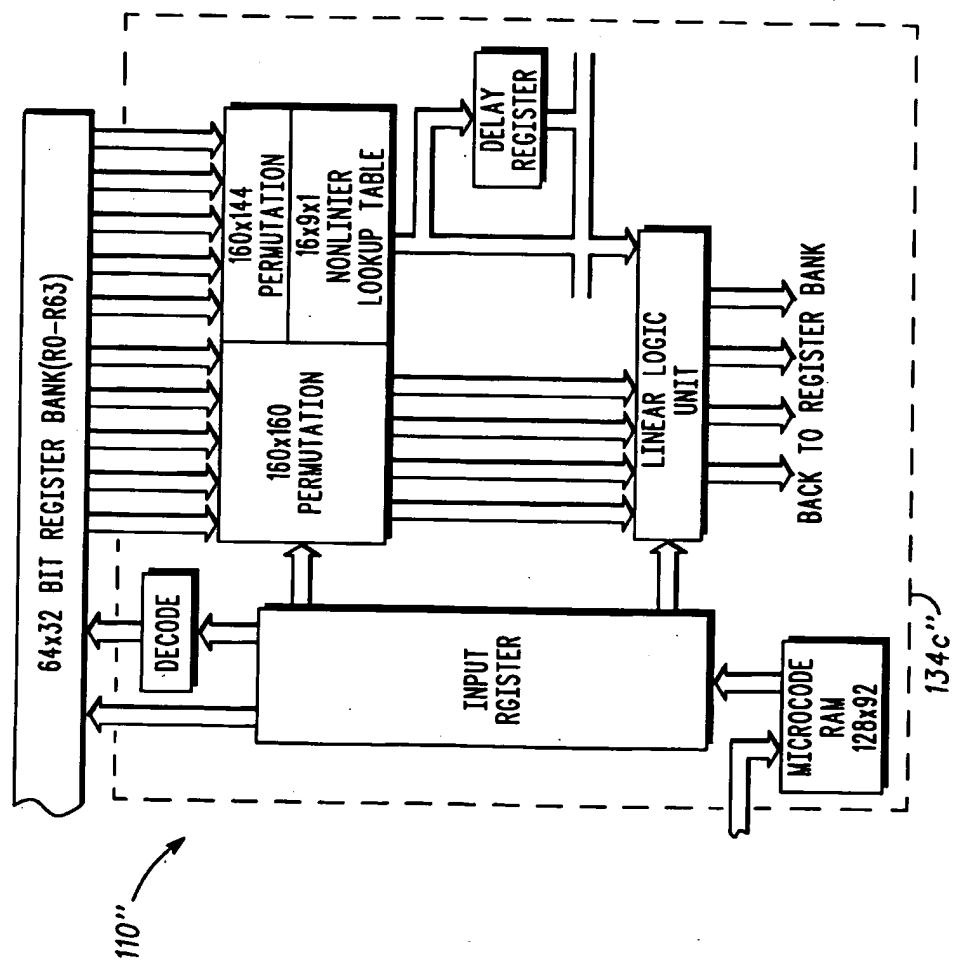


FIG. 11A

FIG. 11B



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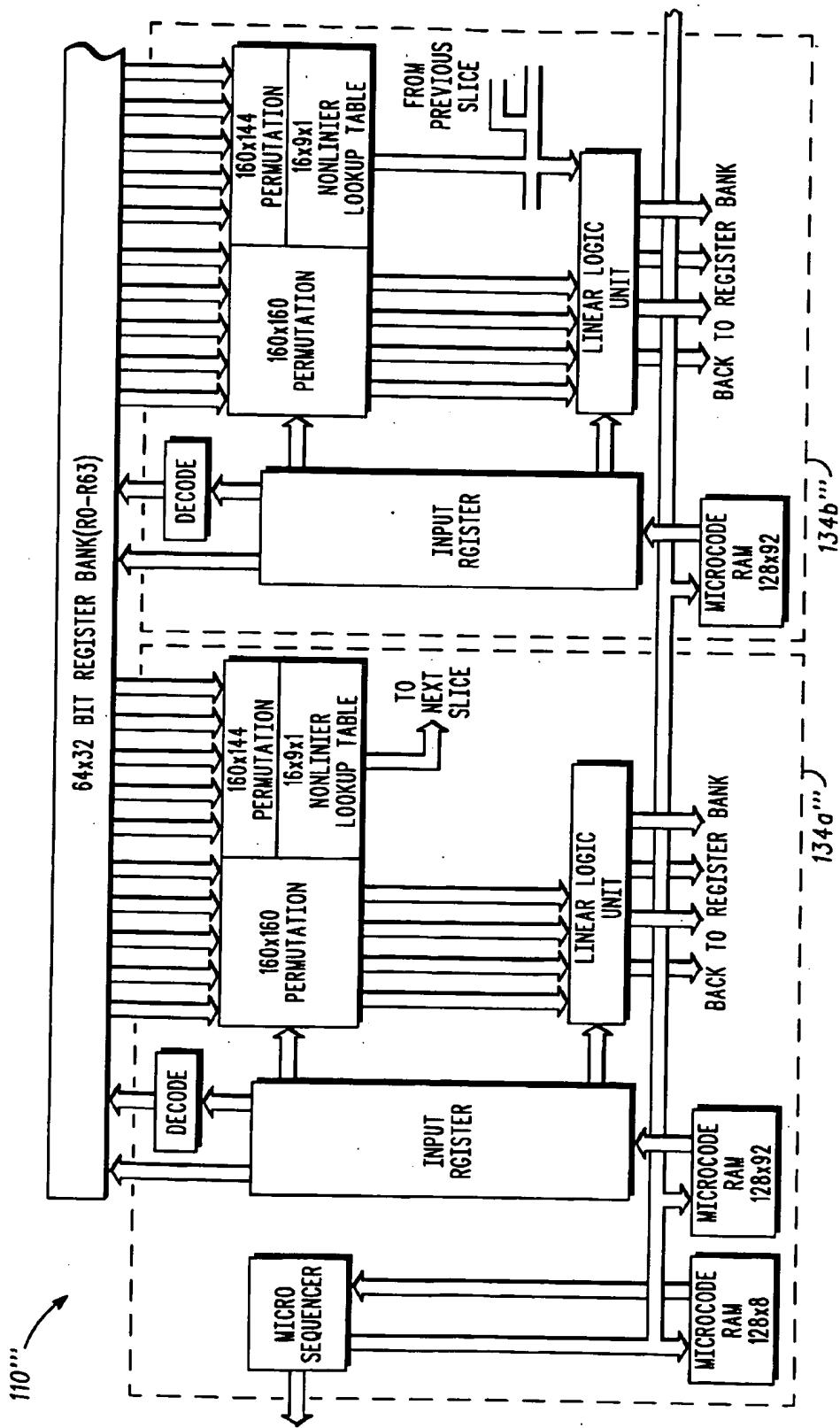


FIG. 12A

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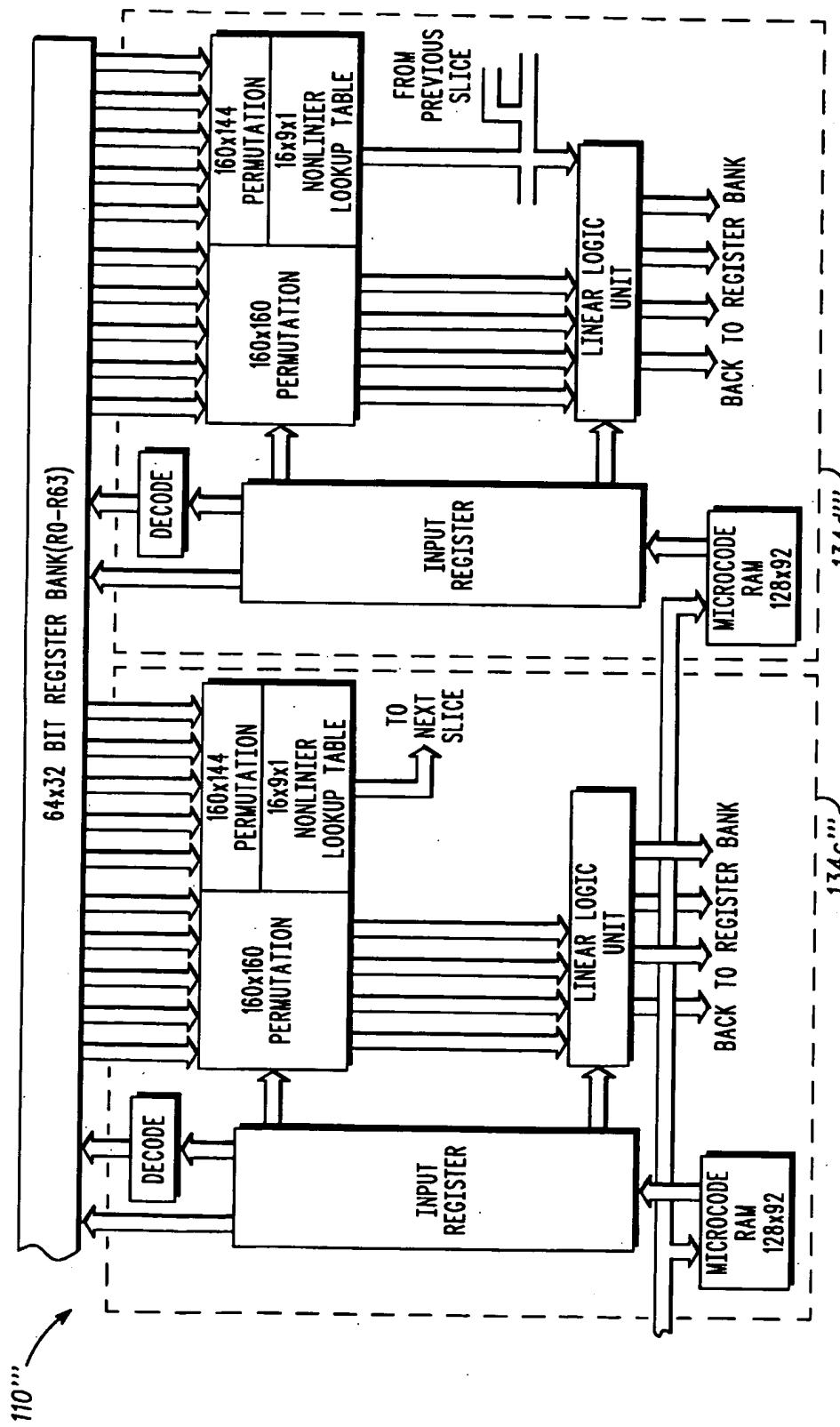


FIG. 12B